## Features

- Single 2.7V - 3.6V Supply
- Serial Peripheral Interface (SPI) Compatible
- Supports SPI Modes 0 and 3
- 66 MHz Maximum Clock Frequency
- Flexible, Uniform Erase Architecture
- 4-Kbyte Blocks
- 32-Kbyte Blocks
- 64-Kbyte Blocks
- Full Chip Erase
- Individual Sector Protection with Global Protect/Unprotect Feature
- Sixty-Four 64-Kbyte Physical Sectors
- Hardware Controlled Locking of Protected Sectors
- Flexible Programming
- Byte/Page Program (1 to 256 Bytes)
- Automatic Checking and Reporting of Erase/Program Failures
- JEDEC Standard Manufacturer and Device ID Read Methodology
- Low Power Dissipation
- 7 mA Active Read Current (Typical)
- $4 \mu \mathrm{~A}$ Deep Power-Down Current (Typical)
- Endurance: 100,000 Program/Erase Cycles
- Data Retention: 20 Years
- Complies with Full Industrial Temperature Range
- Industry Standard Green (Pb/Halide-free/RoHS Compliant) Package Options
- 8-lead SOIC (208-mil wide)
- 16-lead SOIC (300-mil wide)


## 1. Description

The AT26DF321 is a serial interface Flash memory device designed for use in a wide variety of high-volume consumer based applications in which program code is shadowed from Flash memory into embedded or external RAM for execution. The flexible erase architecture of the AT26DF321, with its erase granularity as small as 4-Kbytes, makes it ideal for data storage as well, eliminating the need for additional data storage EEPROM devices.

The physical sectoring and the erase block sizes of the AT26DF321 have been optimized to meet the needs of today's code and data storage applications. By optimizing the size of the physical sectors and erase blocks, the memory space can be used much more efficiently. Because certain code modules and data storage segments must reside by themselves in their own protected sectors, the wasted and unused memory space that occurs with large sectored and large block erase Flash memory devices can be greatly reduced. This increased memory space efficiency allows additional code routines and data storage segments to be added while still maintaining the same overall device density.


32-megabit
2.7-volt Only

Serial Firmware DataFlash ${ }^{\circledR}$ Memory

The AT26DF321 also offers a sophisticated method for protecting individual sectors against erroneous or malicious program and erase operations. By providing the ability to individually protect and unprotect sectors, a system can unprotect a specific sector to modify its contents while keeping the remaining sectors of the memory array securely protected. This is useful in applications where program code is patched or updated on a subroutine or module basis, or in applications where data storage segments need to be modified without running the risk of errant modifications to the program code segments. In addition to individual sector protection capabilities, the AT26DF321 incorporates Global Protect and Global Unprotect features that allow the entire memory array to be either protected or unprotected all at once. This reduces overhead during the manufacturing process since sectors do not have to be unprotected one-by-one prior to initial programming.

Specifically designed for use in 3-volt systems, the AT26DF321 supports read, program, and erase operations with a supply voltage range of 2.7 V to 3.6 V . No separate voltage is required for programming and erasing.

## 2. Pin Descriptions and Pinouts

Table 2-1. Pin Descriptions

| Symbol | Name and Function | Asserted State | Type |
| :---: | :---: | :---: | :---: |
| $\overline{\mathrm{CS}}$ | CHIP SELECT: Asserting the $\overline{\mathrm{CS}}$ pin selects the device. When the $\overline{\mathrm{CS}}$ pin is deasserted, the device will be deselected and normally be placed in standby mode (not Deep Power-Down mode), and the SO pin will be in a high-impedance state. When the device is deselected, data will not be accepted on the SI pin. <br> A high-to-low transition on the $\overline{\mathrm{CS}}$ pin is required to start an operation, and a low-to-high transition is required to end an operation. When ending an internally self-timed operation such as a program or erase cycle, the device will not enter the standby mode until the completion of the operation. | Low | Input |
| SCK | SERIAL CLOCK: This pin is used to provide a clock to the device and is used to control the flow of data to and from the device. Command, address, and input data present on the SI pin is always latched on the rising edge of SCK, while output data on the SO pin is always clocked out on the falling edge of SCK. |  | Input |
| SI | SERIAL INPUT: The SI pin is used to shift data into the device. The SI pin is used for all data input including command and address sequences. Data on the SI pin is always latched on the rising edge of SCK. |  | Input |
| SO | SERIAL OUTPUT: The SO pin is used to shift data out from the device. Data on the SO pin is always clocked out on the falling edge of SCK. |  | Output |
| $\overline{W P}$ | WRITE PROTECT: The $\overline{W P}$ pin controls the hardware locking feature of the device. Please refer to "Protection Commands and Features" on page 11 for more details on protection features and the $\overline{\mathrm{WP}}$ pin. <br> The $\overline{W P}$ pin is internally pulled-high and may be left floating if hardware controlled protection will not be used. However, it is recommended that the $\overline{\mathrm{WP}}$ pin also be externally connected to $\mathrm{V}_{\mathrm{CC}}$ whenever possible. | Low | Input |
| $\mathrm{V}_{\mathrm{Cc}}$ | DEVICE POWER SUPPLY: The $\mathrm{V}_{C C}$ pin is used to supply the source voltage to the device. Operations at invalid $\mathrm{V}_{\mathrm{CC}}$ voltages may produce spurious results and should not be attempted. |  | Power |
| GND | GROUND: The ground reference for the power supply. GND should be connected to the system ground. |  | Power |

Figure 2-1. $\quad$ 8-SOIC Top View


Figure 2-2. 16-SOIC Top View


## 3. Block Diagram



## 4. Memory Array

To provide the greatest flexibility, the memory array of the AT26DF321 can be erased in four levels of granularity including a full chip erase. In addition, the array has been divided into physical sectors of uniform size, of which each sector can be individually protected from program and erase operations. The size of the physical sectors is optimized for both code and data storage applications, allowing both code and data segments to reside in their own isolated regions. Figure $4-1$ on page 4 illustrates the breakdown of each erase level as well as the breakdown of each physical sector.

Figure 4-1. Memory Architecture Diagram

| Internal Sectoring for Sector Protection Function | Block Erase Detail |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | 64KB <br> Block Erase (D8h Command) | 32KB <br> Block Erase <br> (52h Command) | 4KB <br> Block Erase (20h Command) | Block Address Range |
| 64KB <br> (Sector 63) | 64 KB | 32 KB | 4KB | 3FFFFFh - 3FF000h |
|  |  |  | 4KB | 3FEFFFh - 3FE000h |
|  |  |  | 4KB | 3FDFFFh - 3FD000h |
|  |  |  | 4KB | 3FCFFFh - 3FC000h |
|  |  |  | 4KB | 3FBFFFh - 3FB000h |
|  |  |  | 4KB | 3FAFFFh - 3FA000h |
|  |  |  | 4KB | 3F9FFFh - 3F9000h |
|  |  |  | 4KB | 3F8FFFh - 3F8000h |
|  |  | 32 KB | 4KB | $\begin{aligned} & \text { 3F7FFFh - 3F7000h } \\ & \text { 3F6FFFh - 3F6000h } \end{aligned}$ |
|  |  |  | 4KB |  |
|  |  |  | 4KB | 3F5FFFh - 3F5000h <br> 3F4FFFh - 3F4000h |
|  |  |  | 4KB |  |
|  |  |  | 4KB | 3F3FFFh - 3F3000h <br> 3F2FFFh - 3F2000h |
|  |  |  | 4KB |  |
|  |  |  | 4KB | 3F1FFFh - 3F1000h <br> 3FOFFFh - 3F0000h |
|  |  |  | 4KB |  |
| 64KB <br> (Sector 62) | 64 KB | 32 KB | 4KB | 3EFFFFh - 3EF000h |
|  |  |  | 4KB | 3EEFFFh - 3EE000h |
|  |  |  | 4KB | 3EDFFFh - 3ED000h |
|  |  |  | 4KB | 3ECFFFh - 3EC000h |
|  |  |  | 4KB | 3EBFFFh - 3EB000h |
|  |  |  | 4KB | 3EAFFFh - 3EA000h |
|  |  |  | 4KB | 3E9FFFh - 3E9000h |
|  |  |  | 4KB | 3E8FFFh - 3E8000h |
|  |  | 32 KB | 4KB | 3E7FFFh - 3E7000h |
|  |  |  | 4KB | 3E6FFFh - 3E6000h |
|  |  |  | 4KB | 3E5FFFh - 3E5000h |
|  |  |  | 4KB | 3E4FFFh - 3E4000h |
|  |  |  | 4KB | 3E3FFFh - 3E3000h |
|  |  |  | 4KB | 3E2FFFh - 3E2000h |
|  |  |  | 4KB | 3E1FFFh - 3E1000h3E0FFFh - 3E0000h |
|  |  |  | 4KB |  |
| $:$ | : | $:$ | : | 00FFFFh - 00F000h |
| 64KB <br> (Sector 0) | 64 KB | 32 KB | 4KB |  |
|  |  |  | 4KB | 00EFFFh - 00E000h |
|  |  |  | 4KB | 00DFFFh - 00D000h |
|  |  |  | 4KB | 00CFFFh - 00C000h |
|  |  |  | 4KB | 00BFFFh - 00B000h |
|  |  |  | 4KB | 00AFFFh - 00A000h |
|  |  |  | 4KB | 009FFFh - 009000h |
|  |  |  | 4KB | 008FFFh - 008000h |
|  |  | 32 KB | 4KB | 007FFFh - 007000h |
|  |  |  | 4KB | 006FFFh - 006000h |
|  |  |  | 4KB | 005FFFh - 005000h |
|  |  |  | 4KB | 004FFFh - 004000h |
|  |  |  | 4KB | 003FFFh - 003000h |
|  |  |  | 4KB | 002FFFh - 002000h |
|  |  |  | 4KB | 001FFFh - 001000h |
|  |  |  | 4KB | 000FFFh - 000000h |

Page Program Detail

1-256 Byte


| 256 Bytes |
| :--- |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
|  |
| : |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 256 Bytes |
| 25 |
| 2 |
| 2 |

3FFFFFh - 3FFF00h 3FFEFFh - 3FFE00h 3FFDFFh - 3FFD00h 3FFCFFFh - 3FFCOOh 3FFBFFh - 3FFB00h 3FFAFFh - 3FFA00h 3FF9FFh - 3FF900h $3 F F 8 F F h-3 F F 800 h$
$3 F F 7 F F h-3 F F 700 h$ 3FF7FFh - 3FF700h
3FF6FFh - 3FF600h 3FF5FFh - 3FF500h 3FF4FFh - 3FF400h 3FF3FFh - 3FF300h 3FF2FFh - 3FF200h 3FFOFFh - 3FF000h 3FEFFFh - 3FEFOOh 3FEEFFh - 3FEEOOh 3FEDFFh - 3FEDOOh 3FECFFh - 3FEC00h 3FEBFFh - 3FEB00h 3FEAFFh - 3FEA00h 3FE9FFh - 3FE900h
3FE8FFh - 3FE800h

0017FFh - 001700h 0016FFh - 001600h 0015FFh - 001500h 0014FFh - 001400h
0013FFh - 001300h 0012FFh - 001200h 0011FFh - 001100h 0010FFh - 001000h 000FFFh - 000F00h 000EFFh - 000E00h 000DFFh - 000D00h 000CFFh - 000C00h
000BFFh - 000B00h 000AFFh - 000A00h 0009FFh - 000900h 0008FFh - 000800h 0007FFh - 000700h 0006FFh - 000600h 0005FFh - 000500h 0004FFh - 000400h 0003FFh - 000300h 0002FFh - 000200h
0001FFh - 000100h 0000FFh - 000000h

## 5. Device Operation

The AT26DF321 is controlled by a set of instructions that are sent from a host controller, commonly referred to as the SPI Master. The SPI Master communicates with the AT26DF321 via the SPI bus which is comprised of four signal lines: Chip Select ( $\overline{\mathrm{CS}}$ ), Serial Clock (SCK), Serial Input (SI), and Serial Output (SO).

The SPI protocol defines a total of four modes of operation (mode $0,1,2$, or 3 ) with each mode differing in respect to the SCK polarity and phase and how the polarity and phase control the flow of data on the SPI bus. The AT26DF321 supports the two most common modes, SPI Modes 0 and 3 . The only difference between SPI Modes 0 and 3 is the polarity of the SCK signal when in the inactive state (when the SPI Master is in standby mode and not transferring any data). With SPI Modes 0 and 3, data is always latched in on the rising edge of SCK and always output on the falling edge of SCK.

Figure 5-1. SPI Mode 0 and 3


## 6. Commands and Addressing

A valid instruction or operation must always be started by first asserting the $\overline{\mathrm{CS}}$ pin. After the $\overline{\mathrm{CS}}$ pin has been asserted, the SPI Master must then clock out a valid 8 -bit opcode on the SPI bus. Following the opcode, instruction dependent information such as address and data bytes would then be clocked out by the SPI Master. All opcode, address, and data bytes are transferred with the most significant bit (MSB) first. An operation is ended by deasserting the CS pin.

Opcodes not supported by the AT26DF321 will be ignored by the device and no operation will be started. The device will continue to ignore any data presented on the SI pin until the start of the next operation ( $\overline{\mathrm{CS}}$ pin being deasserted and then reasserted). In addition, if the $\overline{\mathrm{CS}}$ pin is deasserted before complete opcode and address information is sent to the device, then no operation will be performed and the device will simply return to the idle state and wait for the next operation.

Addressing of the device requires a total of three bytes of information to be sent, representing address bits A23-A0. Since the upper address limit of the AT26DF321 memory array is 3FFFFFh, address bits A23-A22 are always ignored by the device.

Table 6-1. Command Listing

| Command | Opcode |  | Address Bytes | Dummy Bytes | Data Bytes |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Read Commands |  |  |  |  |  |
| Read Array | 0Bh | 00001011 | 3 | 1 | 1+ |
| Read Array (Low Frequency) | 03h | 00000011 | 3 | 0 | 1+ |
| Program and Erase Commands |  |  |  |  |  |
| Block Erase (4-KBytes) | 20h | 00100000 | 3 | 0 | 0 |
| Block Erase (32-KBytes) | 52h | 01010010 | 3 | 0 | 0 |
| Block Erase (64-KBytes) | D8h | 11011000 | 3 | 0 | 0 |
| Chip Erase | 60h | 01100000 | 0 | 0 | 0 |
|  | C7h | 11000111 | 0 | 0 | 0 |
| Byte/Page Program (1 to 256 Bytes) | 02h | 00000010 | 3 | 0 | 1+ |
| Protection Commands |  |  |  |  |  |
| Write Enable | 06h | 00000110 | 0 | 0 | 0 |
| Write Disable | 04h | 00000100 | 0 | 0 | 0 |
| Protect Sector | 36h | 00110110 | 3 | 0 | 0 |
| Unprotect Sector | 39h | 00111001 | 3 | 0 | 0 |
| Global Protect/Unprotect | Use Write Status Register command |  |  |  |  |
| Read Sector Protection Registers | 3Ch | 00111100 | 3 | 0 | 1+ |
| Status Register Commands |  |  |  |  |  |
| Read Status Register | 05h | 00000101 | 0 | 0 | 1+ |
| Write Status Register | 01h | 00000001 | 0 | 0 | 1 |
| Miscellaneous Commands |  |  |  |  |  |
| Read Manufacturer and Device ID | 9Fh | 10011111 | 0 | 0 | 1 to 4 |
| Deep Power-Down | B9h | 10111001 | 0 | 0 | 0 |
| Resume from Deep Power-Down | ABh | 10101011 | 0 | 0 | 0 |

## 7. Read Commands

### 7.1 Read Array

The Read Array command can be used to sequentially read a continuous stream of data from the device by simply providing the SCK signal once the initial starting address has been specified. The device incorporates an internal address counter that automatically increments on every clock cycle.

Two opcodes, 0Bh and 03h, can be used for the Read Array command. The use of each opcode depends on the maximum SCK frequency that will be used to read data from the device. The OBh opcode can be used at any SCK frequency up to the maximum specified by $f_{\text {ScK }}$. The 03h opcode can be used for lower frequency read operations up to the maximum specified by $f_{\text {RDLF }}$.

To perform the Read Array operation, the $\overline{\mathrm{CS}}$ pin must first be asserted and the appropriate opcode (0Bh or 03h) must be clocked into the device. After the opcode has been clocked in, the three address bytes must be clocked in to specify the starting address location of the first byte to read within the memory array. If the OBh opcode is used, then one don't care byte must also be clocked in after the three address bytes.
After the three address bytes (and the one don't care byte if using opcode 0Bh) have been clocked in, additional clock cycles will result in serial data being output on the SO pin. The data is always output with the MSB of a byte first. When the last byte (3FFFFFh) of the memory array has been read, the device will continue reading back at the beginning of the array (000000h). No delays will be incurred when wrapping around from the end of the array to the beginning of the array.

Deasserting the $\overline{\mathrm{CS}}$ pin will terminate the read operation and put the SO pin into a high-impedance state. The $\overline{\mathrm{CS}}$ pin can be deasserted at any time and does not require that a full byte of data be read.

Figure 7-1. Read Array - OBh Opcode


Figure 7-2. Read Array - 03h Opcode


## 8. Program and Erase Commands

### 8.1 Byte/Page Program

The Byte/Page Program command allows anywhere from a single byte of data to 256 bytes of data to be programmed into previously erased memory locations. An erased memory location is one that has all eight bits set to the logical "1" state (a byte value of FFh). Before a Byte/Page Program command can be started, the Write Enable command must have been previously issued to the device (see Write Enable command description) to set the Write Enable Latch (WEL) bit of the Status Register to a logical "1" state.

To perform a Byte/Page Program command, an opcode of 02h must be clocked into the device followed by the three address bytes denoting the first byte location of the memory array to begin programming at. After the address bytes have been clocked in, data can then be clocked into the device and will be stored in an internal buffer.

If the starting memory address denoted by A23-A0 does not fall on an even 256-byte page boundary (A7-A0 are not all 0), then special circumstances regarding which memory locations will be programmed will apply. In this situation, any data that is sent to the device that goes beyond the end of the page will wrap around back to the beginning of the same page. For example, if the starting address denoted by A23-A0 is 0000FEh, and three bytes of data are sent to the device, then the first two bytes of data will be programmed at addresses 0000FEh and 0000 FFh while the last byte of data will be programmed at address 000000h. The remaining bytes in the page (addresses 000001h through 0000FDh) will be unaffected and will not change. In addition, if more than 256 bytes of data are sent to the device, then only the last 256 bytes sent will be latched into the internal buffer.

When the $\overline{\mathrm{CS}}$ pin is deasserted, the device will take the data stored in the internal buffer and program it into the appropriate memory array locations based on the starting address specified by A23-A0 and the number of complete data bytes sent to the device. If less than 256 bytes of data were sent to the device, then the remaining bytes within the page will not be altered. The programming of the data bytes is internally self-timed and should take place in a time of $t_{p p}$.

The three address bytes and at least one complete byte of data must be clocked into the device before the $\overline{\mathrm{CS}}$ pin is deasserted; otherwise, the device will abort the operation and no data will be programmed into the memory array. In addition, if the address specified by A23-A0 points to a memory location within a sector that is in the protected state (see "Protect Sector" on page 12), then the Byte/Page Program command will not be executed, and the device will return to the idle state once the CS pin has been deasserted. The WEL bit in the Status Register will be reset back to the logical " 0 " state if the program cycle aborts due to an incomplete address being sent, an incomplete byte of data being sent, or because the memory location to be programmed is protected.

While the device is programming, the Status Register can be read and will indicate that the device is busy. For faster throughput, it is recommended that the Status Register be polled rather than waiting the $t_{\text {PP }}$ time to determine if the data bytes have finished programming. At some point before the program cycle completes, the WEL bit in the Status Register will be reset back to the logical "0" state.

The device also incorporates an intelligent programming algorithm that can detect when a byte location fails to program properly. If a programming error arises, it will be indicated by the EPE bit in the Status Register.

Figure 8-1. Byte Program


Figure 8-2. Page Program


### 8.2 Block Erase

A block of $4 \mathrm{~K}-$, 32 K -, or 64 K -bytes can be erased (all bits set to the logical "1" state) in a single operation by using one of three different opcodes for the Block Erase command. An opcode of 20h is used for a 4 K -byte erase, an opcode of 52 h is used for a 32 K -byte erase, and an opcode of D8h is used for a 64K-byte erase. Before a Block Erase command can be started, the Write Enable command must have been previously issued to the device to set the WEL bit of the Status Register to a logical "1" state.

To perform a Block Erase, the $\overline{\mathrm{CS}}$ pin must first be asserted and the appropriate opcode (20h, 52 h , or D8h) must be clocked into the device. After the opcode has been clocked in, the three address bytes specifying an address within the $4 \mathrm{~K}-$, 32 K -, or 64 K -byte block to be erased must be clocked in. Any additional data clocked into the device will be ignored. When the $\overline{\mathrm{CS}}$ pin is deasserted, the device will erase the appropriate block. The erasing of the block is internally self-timed and should take place in a time of $t_{\text {BLKE }}$.

Since the Block Erase command erases a region of bytes, the lower order address bits do not need to be decoded by the device. Therefore, for a 4K-byte erase, address bits A11-A0 will be ignored by the device and their values can be either a logical "1" or "0". For a 32K-byte erase, address bits A14-A0 will be ignored, and for a 64K-byte erase, address bits A15-A0 will be ignored by the device. Despite the lower order address bits not being decoded by the device, the complete three address bytes must still be clocked into the device before the $\overline{\mathrm{CS}}$ pin is deasserted; otherwise, the device will abort the operation and no erase operation will be performed.

If the address specified by A23-A0 points to a memory location within a sector that is in the protected state, then the Block Erase command will not be executed, and the device will return to the idle state once the $\overline{\mathrm{CS}}$ pin has been deasserted.

The WEL bit in the Status Register will be reset back to the logical "0" state if the erase cycle aborts due to an incomplete address being sent or because a memory location within the region to be erased is protected.

While the device is executing a successful erase cycle, the Status Register can be read and will indicate that the device is busy. For faster throughput, it is recommended that the Status Register be polled rather than waiting the $t_{\text {BLKE }}$ time to determine if the device has finished erasing. At some point before the erase cycle completes, the WEL bit in the Status Register will be reset back to the logical "0" state.

The device also incorporates an intelligent erasing algorithm that can detect when a byte location fails to erase properly. If an erase error arises, it will be indicated by the EPE bit in the Status Register.

Figure 8-3. Block Erase


SO HIGH-IMPEDANCE

### 8.3 Chip Erase

The entire memory array can be erased in a single operation by using the Chip Erase command. Before a Chip Erase command can be started, the Write Enable command must have been previously issued to the device to set the WEL bit of the Status Register to a logical "1" state.

Two opcodes, 60h and C7h, can be used for the Chip Erase command. There is no difference in device functionality when utilizing the two opcodes, so they can be used interchangeably. To perform a Chip Erase, one of the two opcodes (60h or C7h) must be clocked into the device. Since the entire memory array is to be erased, no address bytes need to be clocked into the device, and any data clocked in after the opcode will be ignored. When the $\overline{\mathrm{CS}}$ pin is deasserted, the device will erase the entire memory array. The erasing of the device is internally self-timed and should take place in a time of $\mathrm{t}_{\mathrm{CHPE}}$.
The complete opcode must be clocked into the device before the $\overline{C S}$ pin is deasserted; otherwise, no erase will be performed. In addition, if any sector of the memory array is in the protected state, then the Chip Erase command will not be executed, and the device will return to the idle state once the $\overline{\mathrm{CS}}$ pin has been deasserted. The WEL bit in the Status Register will be reset back to the logical " 0 " state if a sector is in the protected state.

While the device is executing a successful erase cycle, the Status Register can be read and will indicate that the device is busy. For faster throughput, it is recommended that the Status Register be polled rather than waiting the $\mathrm{t}_{\mathrm{CHPE}}$ time to determine if the device has finished erasing. At
some point before the erase cycle completes, the WEL bit in the Status Register will be reset back to the logical " 0 " state.

The device also incorporates an intelligent erasing algorithm that can detect when a byte location fails to erase properly. If an erase error arises, it will be indicated by the EPE bit in the Status Register.

Figure 8-4. Chip Erase

$\qquad$

## 9. Protection Commands and Features

### 9.1 Write Enable

The Write Enable command is used to set the Write Enable Latch (WEL) bit in the Status Register to a logical "1" state. The WEL bit must be set before a program, erase, Protect Sector, Unprotect Sector, or Write Status Register command can be executed. This makes the issuance of these commands a two step process, thereby reducing the chances of a command being accidentally or erroneously executed. If the WEL bit in the Status Register is not set prior to the issuance of one of these commands, then the command will not be executed.

To issue the Write Enable command, the $\overline{\mathrm{CS}}$ pin must first be asserted and the opcode of 06h must be clocked into the device. No address bytes need to be clocked into the device, and any data clocked in after the opcode will be ignored. When the $\overline{\mathrm{CS}}$ pin is deasserted, the WEL bit in the Status Register will be set to a logical "1". The complete opcode must be clocked into the device before the $\overline{\mathrm{CS}}$ pin is deasserted; otherwise, the device will abort the operation and the state of the WEL bit will not change.

Figure 9-1. Write Enable


### 9.2 Write Disable

The Write Disable command is used to reset the Write Enable Latch (WEL) bit in the Status Register to the logical "0" state. With the WEL bit reset, all program, erase, Protect Sector, Unprotect Sector, and Write Status Register commands will not be executed. The Write Disable command is also used to exit the Sequential Program Mode. Other conditions can also cause the WEL bit to be reset; for more details, refer to the WEL bit section of the Status Register description on page 20.

To issue the Write Disable command, the $\overline{\mathrm{CS}}$ pin must first be asserted and the opcode of 04 h must be clocked into the device. No address bytes need to be clocked into the device, and any data clocked in after the opcode will be ignored. When the $\overline{C S}$ pin is deasserted, the WEL bit in the Status Register will be reset to a logical " 0 ". The complete opcode must be clocked into the device before the $\overline{\mathrm{CS}}$ pin is deasserted; otherwise, the device will abort the operation and the state of the WEL bit will not change.

Figure 9-2. Write Disable

$\qquad$

### 9.3 Protect Sector

Every physical sector of the device has a corresponding single-bit Sector Protection Register that is used to control the software protection of a sector. Upon device power-up or after a device reset, each Sector Protection Register will default to the logical "1" state indicating that all sectors are protected and cannot be programmed or erased.

Issuing the Protect Sector command to a particular sector address will set the corresponding Sector Protection Register to the logical "1" state. The following table outlines the two states of the Sector Protection Registers.

Table 9-1. $\quad$ Sector Protection Register Values

| Value | Sector Protection Status |
| :---: | :--- |
| 0 | Sector is unprotected and can be programmed and erased. |
| 1 | Sector is protected and cannot be programmed or erased. This is the default state. |

Before the Protect Sector command can be issued, the Write Enable command must have been previously issued to set the WEL bit in the Status Register to a logical "1". To issue the Protect Sector command, the $\overline{\mathrm{CS}}$ pin must first be asserted and the opcode of 36 h must be clocked into the device followed by three address bytes designating any address within the sector to be locked. Any additional data clocked into the device will be ignored. When the CS pin is deasserted, the Sector Protection Register corresponding to the physical sector addressed by A23AO will be set to the logical "1" state, and the sector itself will then be protected from program
and erase operations. In addition, the WEL bit in the Status Register will be reset back to the logical "0" state.

The complete three address bytes must be clocked into the device before the $\overline{\mathrm{CS}}$ pin is deasserted; otherwise, the device will abort the operation, the state of the Sector Protection Register will be unchanged, and the WEL bit in the Status Register will be reset to a logical " 0 ".

As a safeguard against accidental or erroneous protecting or unprotecting of sectors, the Sector Protection Registers can themselves be locked from updates by using the SPRL (Sector Protection Registers Locked) bit of the Status Register (please refer to "Status Register Commands" on page 19 for more details). If the Sector Protection Registers are locked, then any attempts to issue the Protect Sector command will be ignored, and the device will reset the WEL bit in the Status Register back to a logical "0" and return to the idle state once the $\overline{\mathrm{CS}}$ pin has been deasserted.

Figure 9-3. Protect Sector


SO HIGH-IMPEDANCE

### 9.4 Unprotect Sector

Issuing the Unprotect Sector command to a particular sector address will reset the corresponding Sector Protection Register to the logical "0" state (see Table 9-1 on page 12 for Sector Protection Register values). Every physical sector of the device has a corresponding single-bit Sector Protection Register that is used to control the software protection of a sector.

Before the Unprotect Sector command can be issued, the Write Enable command must have been previously issued to set the WEL bit in the Status Register to a logical "1". To issue the Unprotect Sector command, the $\overline{\mathrm{CS}}$ pin must first be asserted and the opcode of 39 h must be clocked into the device. After the opcode has been clocked in, the three address bytes designating any address within the sector to be unlocked must be clocked in. Any additional data clocked into the device after the address bytes will be ignored. When the $\overline{\mathrm{CS}}$ pin is deasserted, the Sector Protection Register corresponding to the sector addressed by A23-A0 will be reset to the logical " 0 " state, and the sector itself will be unprotected. In addition, the WEL bit in the Status Register will be reset back to the logical " 0 " state.
The complete three address bytes must be clocked into the device before the $\overline{\mathrm{CS}}$ pin is deasserted; otherwise, the device will abort the operation, the state of the Sector Protection Register will be unchanged, and the WEL bit in the Status Register will be reset to a logical "0".

As a safeguard against accidental or erroneous locking or unlocking of sectors, the Sector Protection Registers can themselves be locked from updates by using the SPRL (Sector Protection Registers Locked) bit of the Status Register (please refer to "Status Register Commands" on page 19 for more details). If the Sector Protection Registers are locked, then any attempts to
issue the Unprotect Sector command will be ignored, and the device will reset the WEL bit in the Status Register back to a logical "0" and return to the idle state once the $\overline{\mathrm{CS}}$ pin has been deasserted.

Figure 9-4. Unprotect Sector


### 9.5 Global Protect/Unprotect

The Global Protect and Global Unprotect features can work in conjunction with the Protect Sector and Unprotect Sector functions. For example, a system can globally protect the entire memory array and then use the Unprotect Sector command to individually unprotect certain sectors and individually reprotect them later by using the Protect Sector command. Likewise, a system can globally unprotect the entire memory array and then individually protect certain sectors as needed.

Performing a Global Protect or Global Unprotect is accomplished by writing a certain combination of data to the Status Register using the Write Status Register command (see "Write Status Register" section on page 21 for command execution details). The Write Status Register command is also used to modify the SPRL (Sector Protection Registers Locked) bit to control hardware and software locking.
To perform a Global Protect, the appropriate $\overline{W P}$ pin and SPRL conditions must be met and the system must write a logical " 1 " to bits 5,4 , 3 , and 2 of the Status Register. Conversely, to perform a Global Unprotect, the same $\overline{\mathrm{WP}}$ and SPRL conditions must be met but the system must write a logical " 0 " to bits $5,4,3$, and 2 of the Status Register. Table 9-2 details the conditions necessary for a Global Protect or Global Unprotect to be performed.

Table 9-2. Valid SPRL and Global Protect/Unprotect Conditions

| $\begin{aligned} & \overline{\text { WP }} \\ & \text { State } \end{aligned}$ | Current SPRL Value | New Write Status Register Data Bit 76543210 | Protection Operation | New SPRL Value |
| :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  |
| 0 | 0 |  | Global Unprotect - all Sector Protection Registers reset to 0 No change to current protection. <br> No change to current protection. <br> No change to current protection. <br> Global Protect - all Sector Protection Registers set to 1 <br> Global Unprotect - all Sector Protection Registers reset to 0 <br> No change to current protection. <br> No change to current protection. <br> No change to current protection. <br> Global Protect - all Sector Protection Registers set to 1 | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \\ & \\ & 1 \\ & 1 \\ & 1 \\ & 1 \\ & 1 \end{aligned}$ |
| 0 | 1 | x xxxxxx | No change to the current protection level. All sectors currently protected will remain protected and all sectors currently unpro will remain unprotected. <br> The Sector Protection Registers are hard-locked and cannot changed when the $\overline{\mathrm{WP}}$ pin is LOW and the current state of SP Therefore, a Global Protect/Unprotect will not occur. In additio SPRL bit cannot be changed (the WP pin must be HIGH in or change SPRL back to a 0 ). | cted <br> is 1. <br> , the <br> to |
| 1 | 0 |  | Global Unprotect - all Sector Protection Registers reset to 0 No change to current protection. <br> No change to current protection. <br> No change to current protection. <br> Global Protect - all Sector Protection Registers set to 1 <br> Global Unprotect - all Sector Protection Registers reset to 0 <br> No change to current protection. <br> No change to current protection. <br> No change to current protection. <br> Global Protect - all Sector Protection Registers set to 1 | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ |
| 1 | 1 |  | No change to the current protection level. All sectors currently protected will remain protected, and all sectors currently unprotected will remain unprotected. <br> The Sector Protection Registers are soft-locked and cannot be changed when the current state of SPRL is 1 . Therefore, a Global Protect/Unprotect will not occur. However, the SPRL bit can be changed back to a 0 from a 1 since the $\overline{W P}$ pin is HIGH. To perform a Global Protect/Unprotect, the Write Status Register command must be issued again after the SPRL bit has been changed from a 1 to a 0 . | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & 0 \\ & 0 \end{aligned}$ |

Essentially, if the SPRL bit of the Status Register is in the logical " 0 " state (Sector Protection Registers are not locked), then writing a 00 h to the Status Register will perform a Global Unprotect without changing the state of the SPRL bit. Similarly, writing a 7Fh to the Status Register will perform a Global Protect and keep the SPRL bit in the logical "0" state. The SPRL bit can, of course, be changed to a logical "1" by writing an FFh if software-locking or hardware-locking is desired along with the Global Protect.

If the desire is to only change the SPRL bit without performing a Global Protect or Global Unprotect, then the system can simply write a OFh to the Status Register to change the SPRL bit from a logical " 1 " to a logical " 0 " provided the $\overline{\mathrm{WP}}$ pin is deasserted. Likewise, the system can write an FOh to change the SPRL bit from a logical " 0 " to a logical " 1 " without affecting the current sector protection status (no changes will be made to the Sector Protection Registers).

When writing to the Status Register, bits 5, 4, 3, and 2 will not actually be modified but will be decoded by the device for the purposes of the Global Protect and Global Unprotect functions. Only bit 7, the SPRL bit, will actually be modified. Therefore, when reading the Status Register, bits $5,4,3$, and 2 will not reflect the values written to them but will instead indicate the status of the $\overline{W P}$ pin and the sector protection status. Please refer to the "Read Status Register" section and Table 10-1 on page 19 for details on the Status Register format and what values can be read for bits $5,4,3$, and 2 .

### 9.6 Read Sector Protection Registers

The Sector Protection Registers can be read to determine the current software protection status of each sector. Reading the Sector Protection Registers, however, will not determine the status of the $\overline{\mathrm{WP}} \mathrm{pin}$.

To read the Sector Protection Register for a particular sector, the $\overline{C S}$ pin must first be asserted and the opcode of 3Ch must be clocked in. Once the opcode has been clocked in, three address bytes designating any address within the sector must be clocked in. After the last address byte has been clocked in, the device will begin outputting data on the SO pin during every subsequent clock cycle. The data being output will be a repeating byte of either FFh or 00 h to denote the value of the appropriate Sector Protection Register

Table 9-3. Read Sector Protection Register - Output Data

| Output Data | Sector Protection Register Value |
| :---: | :--- |
| 00 h | Sector Protection Register value is 0 (sector is unprotected). |
| FFh | Sector Protection Register value is 1 (sector is protected). |

Deasserting the $\overline{C S}$ pin will terminate the read operation and put the SO pin into a high-impedance state. The $\overline{C S}$ pin can be deasserted at any time and does not require that a full byte of data be read.

In addition to reading the individual Sector Protection Registers, the Software Protection Status (SWP) bit in the Status Register can be read to determine if all, some, or none of the sectors are software protected (please refer to "Status Register Commands" on page 19 for more details).

Figure 9－5．Read Sector Protection Register


## 9．7 Protected States and the Write Protect（ $\overline{\mathrm{WP}}$ ）Pin

The WP pin is not linked to the memory array itself and has no direct effect on the protection sta－ tus of the memory array．Instead，the WP pin，in conjunction with the SPRL（Sector Protection Registers Locked）bit in the Status Register，is used to control the hardware locking mechanism of the device．For hardware locking to be active，two conditions must be met－the $\overline{\mathrm{WP}}$ pin must be asserted and the SPRL bit must be in the logical＂ 1 ＂state．

When hardware locking is active，the Sector Protection Registers are locked and the SPRL bit itself is also locked．Therefore，sectors that are protected will be locked in the protected state， and sectors that are unprotected will be locked in the unprotected state．These states cannot be changed as long as hardware locking is active，so the Protect Sector，Unprotect Sector，and Write Status Register commands will be ignored．In order to modify the protection status of a sector，the $\overline{W P}$ pin must first be deasserted，and the SPRL bit in the Status Register must be reset back to the logical＂0＂state using the Write Status Register command．When resetting the SPRL bit back to a logical＂ 0 ＂，it is not possible to perform a Global Protect or Global Unprotect at the same time since the Sector Protection Registers remain soft－locked until after the Write Status Register command has been executed．

If the $\overline{\mathrm{WP}}$ pin is permanently connected to GND，then once the SPRL bit is set to a logical＂ 1 ＂， the only way to reset the bit back to the logical＂ 0 ＂state is to power－cycle or reset the device． This allows a system to power－up with all sectors software protected but not hardware locked． Therefore，sectors can be unprotected and protected as needed and then hardware locked at a later time by simply setting the SPRL bit in the Status Register．

When the $\overline{W P}$ pin is deasserted，or if the $\overline{W P}$ pin is permanently connected to $V_{C C}$ ，the SPRL bit in the Status Register can still be set to a logical＂ 1 ＂to lock the Sector Protection Registers．This provides a software locking ability to prevent erroneous Protect Sector or Unprotect Sector com－ mands from being processed．When changing the SPRL bit to a logical＂ 1 ＂from a logical＂ 0 ＂，it is also possible to perform a Global Protect or Global Unprotect at the same time by writing the appropriate values into bits 5，4，3，and 2 of the Status Register．

The tables below detail the various protection and locking states of the device.
Table 9-4. Software Protection Register States

| $\mathbf{W P}$ | Sector Protection Register <br> $\mathbf{n}^{(1)}$ | Sector <br> $\mathbf{n}^{(1)}$ |
| :---: | :---: | :---: |
| X | 0 | Unprotected |
| (Don't Care) | 1 | Protected |
|  |  |  |

Note: 1. "n" represents a sector number

Table 9-5. Hardware and Software Locking

| $\overline{\text { WP }}$ | SPRL | Locking | SPRL Change Allowed | Sector Protection Registers |
| :---: | :---: | :--- | :--- | :--- |
| 0 | 0 |  | Can be modified from 0 to 1 | Unlocked and modifiable using the <br> Protect and Unprotect Sector <br> commands. Global Protect and <br> Unprotect can also be performed. |
| 0 | 1 | Hardware <br> Locked | Locked | Locked in current state. Protect and <br> Unprotect Sector commands will be <br> ignored. Global Protect and <br> Unprotect cannot be performed. |
| 1 | 0 |  | Can be modified from 0 to 1 | Unlocked and modifiable using the <br> Protect and Unprotect Sector <br> commands. Global Protect and <br> Unprotect can also be performed. |
| 1 | 1 | Software <br> Locked | Can be modified from 1 to 0 | Locked in current state. Protect and <br> Unprotect Sector commands will be <br> ignored. Global Protect and <br> Unprotect cannot be performed. |

## 10. Status Register Commands

### 10.1 Read Status Register

The Status Register can be read to determine the device's ready/busy status, as well as the status of many other functions such as Hardware Locking and Software Protection. The Status Register can be read at any time, including during an internally self-timed program or erase operation.

To read the Status Register, the $\overline{C S}$ pin must first be asserted and the opcode of 05h must be clocked into the device. After the last bit of the opcode has been clocked in, the device will begin outputting Status Register data on the SO pin during every subsequent clock cycle. After the last bit (bit 0 ) of the Status Register has been clocked out, the sequence will repeat itself starting again with bit 7 as long as the $\overline{\mathrm{CS}}$ pin remains asserted and the SCK pin is being pulsed. The data in the Status Register is constantly being updated, so each repeating sequence will output new data.

Deasserting the $\overline{C S}$ pin will terminate the Read Status Register operation and put the SO pin into a high-impedance state. The $\overline{C S}$ pin can be deasserted at any time and does not require that a full byte of data be read.

Table 10-1. Status Register Format

| Bit ${ }^{(1)}$ |  | Name | Type ${ }^{(2)}$ | Description |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 7 | SPRL | Sector Protection Registers Locked | R/W | $\begin{aligned} & 0 \\ & 1 \end{aligned}$ | Sector Protection Registers are unlocked (default). Sector Protection Registers are locked. |
| 6 | RES | Reserved for future use | R | 0 | Reserved for future use. |
| 5 | EPE | Erase/Program Error | R | $\begin{aligned} & 0 \\ & 1 \end{aligned}$ | Erase or program operation was successful. <br> Erase or program error detected. |
| 4 | WPP | Write Protect ( $\overline{\mathrm{WP}}$ ) Pin Status | R | $\begin{aligned} & 0 \\ & 1 \end{aligned}$ | $\overline{\mathrm{WP}}$ is asserted. <br> $\overline{W P}$ is deasserted. |
| 3:2 | SWP | Software Protection Status | R | 00 <br> 01 <br> 10 <br> 11 | All sectors are software unprotected (all Sector Protection Registers are 0). <br> Some sectors are software protected. Read individual Sector Protection Registers to determine which sectors are protected. <br> Reserved for future use. <br> All sectors are software protected (all Sector Protection Registers are 1 - default). |
| 1 | WEL | Write Enable Latch Status | R | $0$ | Device is not write enabled (default). <br> Device is write enabled. |
| 0 | RDY/BSY | Ready/Busy Status | R | $\begin{aligned} & 0 \\ & 1 \end{aligned}$ | Device is ready. <br> Device is busy with an internal operation. |

Notes: 1. Only bit 7 of the Status Register will be modified when using the Write Status Register command.
2. $R / W=$ Readable and writeable
$R$ = Readable only

### 10.1.1 SPRL Bit

The SPRL bit is used to control whether the Sector Protection Registers can be modified or not. When the SPRL bit is in the logical "1" state, all Sector Protection Registers are locked and cannot be modified with the Protect Sector and Unprotect Sector commands (the device will ignore these commands). In addition, the Global Protect and Global Unprotect features cannot be performed. Any sectors that are presently protected will remain protected, and any sectors that are presently unprotected will remain unprotected.

When the SPRL bit is in the logical " 0 " state, all Sector Protection Registers are unlocked and can be modified (the Protect Sector and Unprotect Sector commands, as well as the Global Protect and Global Unprotect features, will be processed as normal). The SPRL bit defaults to the logical "0" state after a power-up or a device reset.

The SPRL bit can be modified freely whenever the WP pin is deasserted. However, if the WP pin is asserted, then the SPRL bit may only be changed from a logical " 0 " (Sector Protection Registers are unlocked) to a logical "1" (Sector Protection Registers are locked). In order to reset the SPRL bit back to a logical " 0 " using the Write Status Register command, the $\overline{W P}$ pin will have to first be deasserted.

The SPRL bit is the only bit of the Status Register that can be user modified via the Write Status Register command.

### 10.1.2 WPP Bit

### 10.1.3 EPE Bit

The WPP bit can be read to determine if the WP pin has been asserted or not.

The EPE bit indicates whether the last erase or program operation completed successfully or not. If at least one byte during the erase or program operation did not erase or program properly, then the EPE bit will be set to the logical "1" state. The EPE bit will not be set if an erase or program operation aborts for any reason such as an attempt to erase or program a protected region or if the WEL bit is not set prior to an erase or program operation. The EPE bit will be updated after every erase and program operation.

### 10.1.4 SWP Bits

The SWP bits provide feedback on the software protection status for the device. There are three possible combinations of the SWP bits that indicate whether none, some, or all of the sectors have been protected using the Protect Sector command or the Global Protect feature. If the SWP bits indicate that some of the sectors have been protected, then the individual Sector Protection Registers can be read with the Read Sector Protection Registers command to determine which sectors are in fact protected.
10.1.5 WEL Bit

The WEL bit indicates the current status of the internal Write Enable Latch. When the WEL bit is in the logical "0" state, the device will not accept any program, erase, Protect Sector, Unprotect Sector, or Write Status Register commands. The WEL bit defaults to the logical "0" state after a device power-up or reset. In addition, the WEL bit will be reset to the logical "0" state automatically under the following conditions:

- Write Disable operation completes successfully
- Write Status Register operation completes successfully or aborts
- Protect Sector operation completes successfully or aborts
- Unprotect Sector operation completes successfully or aborts
- Byte/Page Program operation completes successfully or aborts
- Block Erase operation completes successfully or aborts
- Chip Erase operation completes successfully or aborts

If the WEL bit is in the logical " 1 " state, it will not be reset to a logical " 0 " if an operation aborts due to an incomplete or unrecognized opcode being clocked into the device before the $\overline{\mathrm{CS}}$ pin is deasserted. In order for the WEL bit to be reset when an operation aborts prematurely, the entire opcode for a program, erase, Protect Sector, Unprotect Sector, or Write Status Register command must have been clocked into the device.

### 10.1.6 RDY/BSY Bit

The RDY/BSY bit is used to determine whether or not an internal operation, such as a program or erase, is in progress. To poll the RDY/BSY bit to detect the completion of a program or erase cycle, new Status Register data must be continually clocked out of the device until the state of the RDY/BSY bit changes from a logical " 1 " to a logical " 0 ".

Figure 10-1. Read Status Register


### 10.2 Write Status Register

The Write Status Register command is used to modify the SPRL bit of the Status Register and/or to perform a Global Protect or Global Unprotect operation. Before the Write Status Register command can be issued, the Write Enable command must have been previously issued to set the WEL bit in the Status Register to a logical "1".

To issue the Write Status Register command, the $\overline{\mathrm{CS}}$ pin must first be asserted and the opcode of 01 h must be clocked into the device followed by one byte of data. The one byte of data consists of the SPRL bit value, a don't care bit, four data bits to denote whether a Global Protect or Unprotect should be performed, and two additional don't care bits (see Table 10-2). Any additional data bytes that are sent to the device will be ignored. When the $\overline{C S}$ pin is deasserted, the SPRL bit in the Status Register will be modified and the WEL bit in the Status Register will be reset back to a logical " 0 ". The values of bits $5,4,3$, and 2 and the state of the SPRL bit before the Write Status Register command was executed (the prior state of the SPRL bit) will determine whether or not a Global Protect or Global Unprotect will be perfomed. Please refer to the "Global Protect/Unprotect" section on page 14 for more details.

The complete one byte of data must be clocked into the device before the $\overline{\mathrm{CS}}$ pin is deasserted; otherwise, the device will abort the operation, the state of the SPRL bit will not change, no potential Global Protect or Unprotect will be performed, and the WEL bit in the Status Register will be reset back to the logical " 0 " state.
If the $\overline{\mathrm{WP}}$ pin is asserted, then the SPRL bit can only be set to a logical " 1 ". If an attempt is made to reset the SPRL bit to a logical "0" while the WP pin is asserted, then the Write Status Register command will be ignored, and the WEL bit in the Status Register will be reset back to the logical " 0 " state. In order to reset the SPRL bit to a logical " 0 ", the WP pin must be deasserted.

Table 10-2. Write Status Register Format

| Bit 7 | Bit 6 | Bit 5 | Bit 4 | Bit 3 | Bit 2 | Bit 1 | Bit 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| SPRL | $X$ | Global Protect/Unprotect |  |  | $X$ | $X$ |  |

Figure 10-2. Write Status Register


## 11. Other Commands and Functions

### 11.1 Read Manufacturer and Device ID

Identification information can be read from the device to enable systems to electronically query and identify the device while it is in system. The identification method and the command opcode comply with the JEDEC standard for "Manufacturer and Device ID Read Methodology for SPI Compatible Serial Interface Memory Devices". The type of information that can be read from the device includes the JEDEC defined Manufacturer ID, the vendor specific Device ID, and the vendor specific Extended Device Information.

To read the identification information, the $\overline{C S}$ pin must first be asserted and the opcode of 9Fh must be clocked into the device. After the opcode has been clocked in, the device will begin outputting the identification data on the SO pin during the subsequent clock cycles. The first byte that will be output will be the Manufacturer ID followed by two bytes of Device ID information. The fourth byte output will be the Extended Device Information String Length, which will be 00h indicating that no Extended Device Information follows. After the Extended Device Information String Length byte is output, the SO pin will go into a high-impedance state; therefore, additional clock cycles will have no affect on the SO pin and no data will be output. As indicated in the JEDEC standard, reading the Extended Device Information String Length and any subsequent data is optional.

Deasserting the $\overline{C S}$ pin will terminate the Manufacturer and Device ID read operation and put the SO pin into a high-impedance state. The $\overline{\mathrm{CS}}$ pin can be deasserted at any time and does not require that a full byte of data be read.

Table 11-1. Manufacturer and Device ID Information

| Byte No. | Data Type | Value |
| :---: | :--- | :---: |
| 1 | Manufacturer ID | 1 h |
| 2 | Device ID (Part 1) | 47 h |
| 3 | Device ID (Part 2) | 00 h |
| 4 | Extended Device Information String Length | 00 h |

Table 11-2. Manufacturer and Device ID Details

| Data Type | Bit 7 | Bit 6 | Bit 5 | Bit 4 | Bit 3 | Bit 2 | Bit 1 | Bit 0 | Hex Value | Details |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Manufacturer ID | JEDEC Assigned Code |  |  |  |  |  |  |  | 1Fh | JEDEC Code: | 00011111 (1Fh for Atmel) |
|  | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 1 |  |  |  |
| Device ID (Part 1) | Family Code |  |  | Density Code |  |  |  |  | 47h | Family Code: Density Code: | 010 (AT25/26DFxxx series) 00111 (32-Mbit) |
|  | 0 | 1 | 0 | 0 | 0 | 1 | 1 | 1 |  |  |  |
| Device ID (Part 2) | Sub Code |  |  | Product Version Code |  |  |  |  | 00h | Sub Code: Product Version: | 000 (Standard series) 00000 (Initial version) |
|  | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |  |  |  |

Figure 11-1. Read Manufacturer and Device ID


### 11.2 Deep Power-Down

During normal operation, the device will be placed in the standby mode to consume less power as long as the CS pin remains deasserted and no internal operation is in progress. The Deep Power-Down command offers the ability to place the device into an even lower power consumption state called the Deep Power-Down mode.

When the device is in the Deep Power-Down mode, all commands including the Read Status Register command will be ignored with the exception of the Resume from Deep Power-Down command. Since all commands will be ignored, the mode can be used as an extra protection mechanism against program and erase operations.

Entering the Deep Power-Down mode is accomplished by simply asserting the $\overline{\mathrm{CS}}$ pin, clocking in the opcode of B9h, and then deasserting the $\overline{\mathrm{CS}}$ pin. Any additional data clocked into the device after the opcode will be ignored. When the $\overline{C S}$ pin is deasserted, the device will enter the Deep Power-Down mode within the maximum time of $\mathrm{t}_{\text {EDPD }}$.
The complete opcode must be clocked in before the $\overline{\mathrm{CS}}$ pin is deasserted; otherwise, the device will abort the operation and return to the standby mode once the $\overline{C S}$ pin is deasserted. In addition, the device will default to the standby mode after a power-cycle or a device reset.

The Deep Power-Down command will be ignored if an internally self-timed operation such as a program or erase cycle is in progress. The Deep Power-Down command must be reissued after the internally self-timed operation has been completed in order for the device to enter the Deep Power-Down mode.

Figure 11-2. Deep Power-Down


### 11.3 Resume from Deep Power-Down

In order exit the Deep Power-Down mode and resume normal device operation, the Resume from Deep Power-Down command must be issued. The Resume from Deep Power-Down command is the only command that the device will recognize while in the Deep Power-Down mode.

To resume from the Deep Power-Down mode, the $\overline{\mathrm{CS}}$ pin must first be asserted and opcode of ABh must be clocked into the device. Any additional data clocked into the device after the opcode will be ignored. When the $\overline{\mathrm{CS}}$ pin is deasserted, the device will exit the Deep PowerDown mode within the maximum time of $t_{\text {RDPD }}$ and return to the standby mode. After the device has returned to the standby mode, normal command operations such as Read Array can be resumed.

If the complete opcode is not clocked in before the $\overline{\mathrm{CS}}$ pin is deasserted, then the device will abort the operation and return to the Deep Power-Down mode.

Figure 11-3. Resume from Deep Power-Down


## 12. Electrical Specifications

### 12.1 Absolute Maximum Ratings*

Temperature Under Bias ............................. $-55^{\circ} \mathrm{C}$ to $+125^{\circ} \mathrm{C}$
Storage Temperature................................... $-65^{\circ} \mathrm{C}$ to $+150^{\circ} \mathrm{C}$
All Input Voltages
(including NC Pins)
with Respect to Ground ...................................-0.6V to +4.1 V
All Output Voltages
with Respect to Ground ............................. 0.6 V to $\mathrm{V}_{\mathrm{CC}}+0.5 \mathrm{~V}$
*NOTICE: Stresses beyond those listed under "Absolute Maximum Ratings" may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions beyond those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods may affect device reliability.

### 12.2 DC and AC Operating Range

|  |  | AT26DF321 |
| :--- | :--- | :---: |
| Operating Temperature (Case) | Ind. | $-40^{\circ} \mathrm{C}$ to $+85^{\circ} \mathrm{C}$ |
| $\mathrm{V}_{\mathrm{CC}}$ Power Supply | 2.7 V to 3.6 V |  |

### 12.3 DC Characteristics

| Symbol | Parameter | Condition | Min | Typ | Max | Units |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{I}_{\text {SB }}$ | Standby Current | $\overline{\mathrm{CS}}, \overline{\mathrm{WP}}=\mathrm{V}_{\mathrm{CC}},$ <br> all inputs at CMOS levels |  | 25 | 35 | $\mu \mathrm{A}$ |
| $\mathrm{I}_{\text {DPD }}$ | Deep Power-Down Current | $\overline{\mathrm{CS}}, \overline{\mathrm{WP}}=\mathrm{VCC},$ <br> all inputs at CMOS levels |  | 4 | 8 | $\mu \mathrm{A}$ |
| $\mathrm{I}_{\mathrm{CC1}}$ | Active Current, Read Operation | $\begin{aligned} & \mathrm{f}=66 \mathrm{MHz}, \mathrm{I}_{\mathrm{OUT}}=0 \mathrm{~mA}, \\ & \overline{\mathrm{CS}}=\mathrm{V}_{\mathrm{IL}}, \mathrm{~V}_{\mathrm{CC}}=\mathrm{Max} \end{aligned}$ |  | 11 | 15 | mA |
|  |  | $\begin{aligned} & \mathrm{f}=50 \mathrm{MHz} ; \mathrm{I}_{\text {OUT }}=0 \mathrm{~mA}, \\ & \overline{\mathrm{CS}}=\mathrm{V}_{\mathrm{IL}}, \mathrm{~V}_{\mathrm{CC}}=\mathrm{Max} \end{aligned}$ |  | 10 | 14 |  |
|  |  | $\begin{aligned} & \mathrm{f}=33 \mathrm{MHz}, \mathrm{I}_{\mathrm{OUT}}=0 \mathrm{~mA}, \\ & \overline{\mathrm{CS}}=\mathrm{V}_{\mathrm{IL}}, \mathrm{~V}_{\mathrm{CC}}=\mathrm{Max} \end{aligned}$ |  | 8 | 12 |  |
|  |  | $\begin{aligned} & \mathrm{f}=20 \mathrm{MHz}, \mathrm{I}_{\mathrm{OUT}}=0 \mathrm{~mA}, \\ & \overline{\mathrm{CS}}=\mathrm{V}_{\mathrm{IL}}, \mathrm{~V}_{\mathrm{CC}}=\mathrm{Max} \end{aligned}$ |  | 7 | 10 |  |
| $\mathrm{I}_{\mathrm{CC} 2}$ | Active Current, Program Operation | $\overline{\mathrm{CS}}=\mathrm{V}_{\mathrm{CC}}, \mathrm{V}_{\mathrm{CC}}=\mathrm{Max}$ |  | 12 | 18 | mA |
| $\mathrm{I}_{\text {CC3 }}$ | Active Current, Erase Operation | $\mathrm{cs}=\mathrm{V}_{\mathrm{CC}}, \mathrm{V}_{\mathrm{CC}}=\mathrm{Max}$ |  | 14 | 20 | mA |
| $\mathrm{I}_{\mathrm{LI}}$ | Input Leakage Current | $\mathrm{V}_{\mathrm{IN}}=$ CMOS levels |  |  | 1 | $\mu \mathrm{A}$ |
| $\mathrm{I}_{\text {LO }}$ | Output Leakage Current | $\mathrm{V}_{\text {OUT }}=$ CMOS levels |  |  | 1 | $\mu \mathrm{A}$ |
| $\mathrm{V}_{\text {IL }}$ | Input Low Voltage |  |  |  | $0.3 \times \mathrm{V}_{\mathrm{cc}}$ | V |
| $\mathrm{V}_{\mathrm{IH}}$ | Input High Voltage |  | $0.7 \times \mathrm{V}_{\text {cc }}$ |  |  | V |
| $\mathrm{V}_{\text {OL }}$ | Output Low Voltage | $\mathrm{I}_{\mathrm{OL}}=1.6 \mathrm{~mA}, \mathrm{~V}_{\mathrm{CC}}=\mathrm{Min}$ |  |  | 0.4 | V |
| $\mathrm{V}_{\mathrm{OH}}$ | Output High Voltage | $\mathrm{I}_{\mathrm{OH}}=-100 \mu \mathrm{~A}$ | $\mathrm{V}_{\mathrm{CC}}-0.2$ |  |  | V |

### 12.4 AC Characteristics

| Symbol | Parameter | Min | Max | Units |
| :---: | :---: | :---: | :---: | :---: |
| $\mathrm{f}_{\text {SCK }}$ | Serial Clock (SCK) Frequency |  | 66 | MHz |
| $\mathrm{f}_{\text {RDLF }}$ | SCK Frequency for Read Array (Low Frequency - 03h opcode) |  | 33 | MHz |
| $\mathrm{t}_{\text {SCKH }}$ | SCK High Time | 6.8 |  | ns |
| $\mathrm{t}_{\text {SCKL }}$ | SCK Low Time | 6.8 |  | ns |
| $\mathrm{t}_{\text {SCKR }}{ }^{(1)}$ | SCK Rise Time, Peak-to-Peak (Slew Rate) | 0.1 |  | V/ns |
| $\mathrm{t}_{\text {SCKF }}{ }^{(1)}$ | SCK Fall Time, Peak-to-Peak (Slew Rate) | 0.1 |  | V/ns |
| $\mathrm{t}_{\text {CSH }}$ | Chip Select High Time | 50 |  | ns |
| $\mathrm{t}_{\text {CSLS }}$ | Chip Select Low Setup Time (relative to SCK) | 5 |  | ns |
| $\mathrm{t}_{\mathrm{CSLH}}$ | Chip Select Low Hold Time (relative to SCK) | 5 |  | ns |
| $\mathrm{t}_{\mathrm{CSHS}}$ | Chip Select High Setup Time (relative to SCK) | 5 |  | ns |
| $\mathrm{t}_{\mathrm{CSHH}}$ | Chip Select High Hold Time (relative to SCK) | 5 |  | ns |
| $\mathrm{t}_{\mathrm{DS}}$ | Data In Setup Time | 2 |  | ns |
| $\mathrm{t}_{\mathrm{DH}}$ | Data In Hold Time | 3 |  | ns |
| $\mathrm{t}_{\text {DIS }}{ }^{(1)}$ | Output Disable Time |  | 6 | ns |
| $\mathrm{t}_{\mathrm{v}}$ | Output Valid Time |  | 6 | ns |
| $\mathrm{t}_{\mathrm{OH}}$ | Output Hold Time | 0 |  | ns |
| $\mathrm{t}_{\text {WPS }}{ }^{(1)(2)}$ | Write Protect Setup Time | 20 |  | ns |
| $\mathrm{t}_{\text {WPH }}{ }^{(1)(2)}$ | Write Protect Hold Time | 100 |  | ns |
| $\mathrm{t}_{\text {SECP }}{ }^{(1)}$ | Sector Protect Time (from Chip Select High) |  | 20 | ns |
| $\mathrm{t}_{\text {SECUP }}{ }^{(1)}$ | Sector Unprotect Time (from Chip Select High) |  | 20 | ns |
| $\mathrm{t}_{\text {EDPD }}{ }^{(1)}$ | Chip Select High to Deep Power-Down |  | 3 | $\mu \mathrm{s}$ |
| $\mathrm{t}_{\text {RDPD }}{ }^{(1)}$ | Chip Select High to Standby Mode |  | 3 | $\mu \mathrm{s}$ |

Notes: 1. Not $100 \%$ tested (value guaranteed by design and characterization).
2. Only applicable as a constraint for the Write Status Register command when SPRL = 1 .

### 12.5 Program and Erase Characteristics

| Symbol | Parameter |  | Min | Typ | Max | Units |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{t}_{\mathrm{pP}}$ | Page Program Time (256 Bytes) |  |  | 1.5 | 5.0 | ms |
| $\mathrm{t}_{\mathrm{BP}}$ | Byte Program Time |  |  | 6 |  | $\mu \mathrm{s}$ |
| $t_{\text {BLKE }}$ | Block Erase Time | 4-Kbyte |  | 50 | 200 | ms |
|  |  | 32-Kbyte |  | 350 | 600 |  |
|  |  | 64-Kbyte |  | 600 | 950 |  |
| $\mathrm{t}_{\text {CHPE }}{ }^{(1)}$ | Chip Erase Time |  |  | 36 | 56 | sec |
| $\mathrm{t}_{\text {WRSR }}{ }^{(1)}$ | Write Status Register Time |  |  |  | 200 | ns |

Notes: 1. Not $100 \%$ tested (value guaranteed by design and characterization).

## AIIIE

### 12.6 Power-Up Conditions

| Parameter | Min | Max |
| :--- | :---: | :---: |
| Minimum V ${ }_{\text {cc }}$ to Chip Select Low Time | 50 |  |
| Power-up Device Delay Before Program or Erase Allowed |  | $\mu \mathrm{s}$ |
| Power-On Reset Voltage | 1.5 | 2.5 |

### 12.7 Input Test Waveforms and Measurement Levels



### 12.8 Output Test Load



## 13. AC Waveforms

Figure 13-1. Serial Input Timing


Figure 13-2. Serial Output Timing


Figure 13-3. $\overline{\mathrm{WP}}$ Timing for Write Status Register Command When $\mathrm{SPRL}=1$


## 14. Ordering Information

### 14.1 Green Package Options (Pb/Halide-free/RoHS Compliant)

| $\mathbf{f}_{\mathbf{S C K}}(\mathbf{M H z})$ | Ordering Code | Package | Operation Range |
| :---: | :---: | :---: | :---: |
| 66 | AT26DF321-SU | 8 S 2 | Industrial |
| 66 | AT26DF321-S3U | 16 S | $\left(-40^{\circ} \mathrm{C}\right.$ to $\left.+85^{\circ} \mathrm{C}\right)$ |


| Package Type |  |
| :--- | :--- |
| 8S2 | 8-lead, 0.208" Wide, Plastic Gull Wing Small Outline Package (EIAJ SOIC) |
| $16 S$ | 16-lead, 0.300" Wide, Plastic Gull Wing Small Outline Package (SOIC) |

## 15. Packaging Information

## $15.1 \quad 8 \mathrm{~S} 2$ - EIAJ SOIC




END VIEW


SIDE VIEW

| COMMON DIMENSIONS <br> (Unit of Measure $=\mathrm{mm}$ ) |  |  |  |  |
| :--- | :---: | :---: | :---: | :---: |
| SYMBOL MIN NOM MAX NOTE <br> A 1.70  2.16  <br> A1 0.05  0.25  <br> b 0.35  0.48 4 <br> C 0.15  0.35 4 <br> D 5.13  5.35  <br> E1 5.18  5.40 2 <br> E 7.70  8.26  <br> L 0.51  0.85  <br> $\theta$ $0^{\circ}$  $8^{\circ}$  <br> e 1.27 BSC    |  |  |  |  |

Notes: 1. This drawing is for general information only; refer to EIAJ Drawing EDR-7320 for additional information.
2. Mismatch of the upper and lower dies and resin burrs aren't included.
3. Determines the true geometric position.
4. Values $b, C$ apply to plated terminal. The standard thickness of the plating layer shall measure between 0.007 to .021 mm .

4/15/08

| Package Drawing Contact: packagedrawings@atmel.com | TITLE <br> 8S2, 8-lead, 0.208" Body, Plastic Small Outline Package (EIAJ) | $\begin{aligned} & \hline \text { GPC } \\ & \text { STN } \end{aligned}$ | DRAWING NO. 8S2 | REV. <br> F |
| :---: | :---: | :---: | :---: | :---: |

15.2 16S－SOIC


Notes：1．This drawing is for general information only；refer to JEDEC Drawing MS－013，Variation AA for additional information．
2．Dimension D does not include mold Flash，protrusions or gate burrs．Mold Flash，protrusion and gate burrs shall not exceed $0.15 \mathrm{~mm}\left(0.006{ }^{\prime \prime}\right)$ per side．
3．Dimension E does not include inter－lead Flash or protrusion．Inter－lead flash and protrusions shall not exceed 0.25 mm （0．010＂）per side．
4．$L$ is the length of the terminal for soldering to a substrate．

## 16. Revision History

| Revision Level - Release Date | History |
| :---: | :---: |
| A - May 2006 | Initial release |
| B - July 2006 | Corrected typographical errors. |
| C - August 2006 | Changed description of bit 5 in Table 10-1 from " 0 " to " $x$ " and specified that the value is undefined. <br> Removed MLF package offerings. <br> Added 8-lead SOIC (200-mil wide) package. <br> Changed ordering code for 16-lead SOIC from AT26DF321-SU to AT26DF321-S3U. <br> Added errata regarding Chip Erase. |
| D - April 2007 | Removed "Preliminary" from datasheet. <br> Added information regarding EPE bit. <br> - Added text to "Byte/Page Program" section. <br> - Added text to "Block Erase" section. <br> - Added text to "Chip Erase" section. <br> - Added EPE bit description (bit 5) to Table 10-1. <br> Improved erase times in "Program and Erase Characteristics" table. <br> - Reduced 64KB typical block erase time from 700 ms to 600 ms . <br> - Reduced 64KB maximum block erase time from 1000 ms to 950 ms . |
| E - October 2007 | Added statement recommending AT25DF321 for new designs. |
| F - May 2008 | Removed Section 17(Errata) from the datasheet. |

## Headquarters

Atmel Corporation
2325 Orchard Parkway
San Jose, CA 95131
USA
Tel: 1(408) 441-0311
Fax: 1(408) 487-2600

International

| Atmel Asia | Atmel Europe |
| :--- | :--- |
| Room 1219 | Le Krebs |
| Chinachem Golden Plaza | 8, Rue Jean-Pierre Timbaud |
| 77 Mody Road Tsimshatsui | BP 309 |
| East Kowloon | 78054 Saint-Quentin-en- |
| Hong Kong | Yvelines Cedex |
| Tel: (852) 2721-9778 | France |
| Fax: (852) 2722-1369 | Tel: (33) 1-30-60-70-00 |
|  | Fax: (33) 1-30-60-71-11 |

## Atmel Japan

9F, Tonetsu Shinkawa Bldg.
1-24-8 Shinkawa
Chuo-ku, Tokyo 104-0033
Japan
Tel: (81) 3-3523-3551
Fax: (81) 3-3523-7581

## Product Contact

## Web Site

www.atmel.com

## Technical Support

dataflash@atmel.com

## Sales Contact

www.atmel.com/contacts

## Literature Requests

www.atmel.com/literature

Disclaimer: The information in this document is provided in connection with Atmel products. No license, express or implied, by estoppel or otherwise, to any intellectual property right is granted by this document or in connection with the sale of Atmel products. EXCEPT AS SET FORTH IN ATMEL'S TERMS AND CONDITIONS OF SALE LOCATED ON ATMEL'S WEB SITE, ATMEL ASSUMES NO LIABILITY WHATSOEVER AND DISCLAIMS ANY EXPRESS, IMPLIED OR STATUTORY WARRANTY RELATING TO ITS PRODUCTS INCLUDING, BUT NOT LIMITED TO, THE IMPLIED WARRANTY OF MERCHANTABILITY, FITNESS FOR A PARTICULAR PURPOSE, OR NON-INFRINGEMENT. IN NO EVENT SHALL ATMEL BE LIABLE FOR ANY DIRECT, INDIRECT, CONSEQUENTIAL, PUNITIVE, SPECIAL OR INCIDENTAL DAMAGES (INCLUDING, WITHOUT LIMITATION, DAMAGES FOR LOSS OF PROFITS, BUSINESS INTERRUPTION, OR LOSS OF INFORMATION) ARISING OUT OF THE USE OR INABILITY TO USE THIS DOCUMENT, EVEN IF ATMEL HAS BEEN ADVISED OF THE POSSIBILITY OF SUCH DAMAGES. Atmel makes no representations or warranties with respect to the accuracy or completeness of the contents of this document and reserves the right to make changes to specifications and product descriptions at any time without notice. Atmel does not make any commitment to update the information contained herein. Unless specifically provided otherwise, Atmel products are not suitable for, and shall not be used in, automotive applications. Atmel's products are not intended, authorized, or warranted for use as components in applications intended to support or sustain life.
© 2008 Atmel Corporation. All rights reserved. Atmel ${ }^{\circledR}$, logo and combinations thereof, Everywhere You Are ${ }^{\circledR}$, DataFlash ${ }^{\circledR}$ and others are registered trademarks or trademarks of Atmel Corporation or its subsidiaries. Other terms and product names may be trademarks of others.

